## ON THE CONVEXITY OF STAMPFLI'S NUMERICAL RANGE

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This paper investigates a certain type of numerical range introduced by Stampfli. In particular, we investigate the convexity of this set of elements of operators on Hilbert spaces and its relationship to the algebra numerical range implemented by elements of a  $W^*$ -algebra.

## 1. Introduction

Let  $\mathfrak S$  be a complex Hilbert space,  $\mathfrak L(\mathfrak S)$  the set of bounded linear operators on  $\mathfrak S$ . Stampfli [4] defined the  $\delta$ -numerical range,  $W_{\delta}(T)$ , as the set

$$W_{\delta}(T) = \operatorname{closure}\{\langle Tx, x \rangle \colon ||x|| = 1 \text{ and } ||Tx|| \geqslant \delta\}.$$

Having introduced this set, he wondered whether for any  $T \in \mathfrak{B}(\mathfrak{X})$ , the set  $W_{\delta}(T)$  is convex. (Here  $\mathfrak{B}(\mathfrak{X})$  is the set of bounded linear operators on a Banach space  $\mathfrak{X}$ .) We shall solve this problem for  $\mathfrak{L}(\mathfrak{S})$  using an idea of Dekker, Bonsall and Duncan, [2]. We then proceed to introduce the set  $V_{\delta}(T)$ , where

$$V_{\delta}(T) = \operatorname{closure}\{f(T): f(I) = ||f|| = 1 \text{ and } f(T^*T) \ge \delta^2\}$$

is the algebra  $\delta$ -numerical range for an element T of a unital  $W^*$ -algebra  $\mathfrak{U}$ . We conclude by proving that the sets  $W_{\delta}(T)$  and  $V_{\delta}(T)$  are actually equal.

We remind the reader that a subset G of  $\mathbb{C}$  is said to be connected if it cannot be split into two nonempty open sets. The following lemma shows when the set is connected.

**LEMMA 1.1.** Let T be a self adjoint element of  $\mathfrak{L}(\mathfrak{S})$  and let

$$E = \{x \in \mathfrak{S} : ||x|| \text{ and } \langle Tx, x \rangle = 0\}.$$

Then E is arcwise connected.

The proof of this lemma can be found in Bonsall and Duncan, [2]. We shall now prove the convexity of the set  $W_{\delta}(T)$  for any  $T \in \mathfrak{L}(\mathfrak{S})$ . We would like to point out that the proof of this theorem also appears in the author's Ph.D. thesis, Agure [1].

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THEOREM 1.2.  $W_{\delta}(T)$  is a convex set.

PROOF: Let l be a straight line in  $\mathbb{C}$ . To prove that the set  $W_{\delta}(T)$  is convex, it is enough to show that the intersection  $W_{\delta}(T) \cap l$  is a connected set. If  $T = T_1 + iT_2$ ,  $T_i = T_i^*$ , i = 1, 2 and  $L = aT_1 + bT_2 + cI$ , (where I is the identity operator) and a, b, c are constants, then it is clear that L is a self adjoint operator. Consider the set E given by

$$E = \{x \in \mathfrak{S} \colon ||x|| = 1, ||(T_1 + iT_2)x|| \geqslant \delta \text{ and } \langle Lx, x \rangle = 0\}.$$

Assume that E is nonempty. We shall show that E is arcwise connected. To do that choose any two vectors  $y_1$ ,  $y_2$  in E. Note that there is no loss of generality in assuming that the vectors  $y_1$  and  $y_2$  are linearly independent elements of E. Clearly if  $y \in E$ , then for any  $\alpha \in \mathbb{R}$ ,  $ye^{i\alpha} \in E$  and y is joined to  $ye^{i\alpha}$  by an arc t that maps  $[0, \alpha]$  to E,  $t \to e^{i\alpha t}y$ . Next choose  $\alpha \in \mathbb{R}$  such that  $e^{i\alpha}\langle Ly_1, y_2\rangle$  is purely imaginary and take  $\widehat{y}_1 = y_1e^{i\alpha}$ ,  $\alpha \in \mathbb{R}$ . Secondly take  $\beta$  so that

(\*) 
$$\Re e \langle (T^*T - \delta^2) \widehat{y}_1, e^{i\beta} y_2 \rangle \geqslant 0$$

and set  $\widehat{y}_2 = e^{i\beta}y_2$ . We are left with the task of showing that  $\widehat{y}_1$  can be joined to  $\widehat{y}_2$  by an arc in E. It is clear that  $\langle L\widehat{y}_1, \widehat{y}_2 \rangle$  will be purely imaginary. Now with  $0 \leq s \leq 1$ , let x(s) and p(s) be given by

$$x(s) = (1-s)\widehat{y}_1 + s\widehat{y}_2, \quad p(s) = \frac{1}{\|x(s)\|}x(s).$$

We shall show that the arc p(s) is in the set E by first proving that  $\langle Lp(s), p(s) \rangle = 0$ , and then  $||Tp(s)|| \ge \delta$ . Now

$$\begin{split} \langle Lx(s),\,x(s)\rangle &= \langle L(1-s)\widehat{y}_1+s\widehat{y}_2,\,(1-s)\widehat{y}_1+s\widehat{y}_2\rangle \\ &= (1-s)^2\langle L\widehat{y}_1,\,\widehat{y}_1\rangle + (1-s)s(\langle L\widehat{y}_1,\,\widehat{y}_2\rangle + \langle L\widehat{y}_2,\,\widehat{y}_1\rangle) + s^2\langle L\widehat{y}_2,\,\widehat{y}_2\rangle \\ &= (1-s)s\{\langle L\widehat{y}_1,\,\widehat{y}_2\rangle + \langle L\widehat{y}_2,\,\widehat{y}_1\rangle\} = 2(1-s)s\Re\left(\langle L\widehat{y}_1,\,\widehat{y}_2\rangle = 0 \end{split}$$

as  $\langle L\widehat{y}_1,\,\widehat{y}_2\rangle$  is purely imaginary. Therefore it is clear that  $\langle Lp(s),\,p(s)\rangle=0$ . Also

$$\begin{split} \left\| Tp(s) \right\|^2 &= \frac{1}{\left\| x(s) \right\|^2} \left\| Tx(s) \right\|^2 = \frac{1}{\left\| x(s) \right\|^2} \langle Tx(s), Tx(s) \rangle \\ &= \frac{1}{\left\| x(s) \right\|^2} \left\{ (1-s)^2 \left\| T\widehat{y}_1 \right\|^2 + s^2 \left\| T\widehat{y}_2 \right\|^2 + 2(1-s)s\delta^2 \Re \left\langle \widehat{y}_1, \widehat{y}_2 \right\rangle \right. \\ &+ 2(1-s)s\Re \left\langle T\widehat{y}_1, T\widehat{y}_2 \right\rangle - 2\delta^2 (1-s)s\Re \left\langle \widehat{y}_1, \widehat{y}_2 \right\rangle \right\} \\ &\geqslant \frac{1}{\left\| x(s) \right\|^2} \left\{ \left\| x(s) \right\|^2 \delta^2 + 2(1-s)s\delta^2 \Re \left\langle \widehat{y}_1, \widehat{y}_2 \right\rangle + \\ &\quad 2(1-s)s\Re \left\langle T\widehat{y}_1, T\widehat{y}_2 \right\rangle - 2\delta^2 (1-s)s\Re \left\langle \widehat{y}_1, \widehat{y}_2 \right\rangle \right\} \\ &= \delta^2 + \frac{1}{\left\| x(s) \right\|^2} \left\{ 2(1-s)s\Re \left\langle \left( T^*T - \delta^2 \right) \widehat{y}_1, \widehat{y}_2 \right\rangle \right\}. \end{split}$$

And we see that  $||Tp(s)|| \ge \delta$  by virtue of (\*). We therefore conclude that the set E is connected and the required arc is p(s). Now define the mapping  $\pi$  on E by

$$\pi(x) = \langle T_1 x, x \rangle + i \langle T_2 x, x \rangle, \ \forall x \in E.$$

Then it is clear that  $\pi(x) \in l$  and  $\pi(x) \in W_{\delta}(T)$ . So  $W_{\delta}(T) \cap l = \pi(x)$ . This mapping is also continuous. Since E is connected and the mapping  $\pi$  is continuous we conclude that  $\pi(E)$  is connected and hence  $W_{\delta}(T)$  is convex.

We shall now introduce the algebra numerical range. Suppose  $\mathfrak U$  is a unital  $W^*$ -algebra,  $\mathfrak U^*$  its dual and  $E(\mathfrak U)$  the set of states on  $\mathfrak U$ , that is

$$E(\mathfrak{U}) = \{ f \in \mathfrak{U}^* : f(I) = 1 = ||f|| \}.$$

Then for any T in  $\mathfrak{U}$  we define  $V_{\delta}(T)$  by

$$V_{\delta}(Y) = \operatorname{closure}\{f(T) : f \in E(\mathfrak{U}) \text{ and } f(T^*T) \geqslant \delta^2\}.$$

THEOREM 1.3.  $V_{\delta}(T)$  is a convex set.

PROOF: Let  $\lambda_1$  and  $\lambda_2 \in V_\delta(T)$ . Then there exists states  $g_1$  and  $g_2$  on  $\mathfrak U$  such that  $g_1(T) = \lambda_1$ ,  $g_1(T^*T) \geqslant \delta^2$  and  $g_2(T) = \lambda_2$ ,  $g_2(T^*T) \geqslant \delta^2$ . Now for  $0 \leqslant \alpha \leqslant 1$ , let

$$g(T) = \alpha g_1(T) + (1 - \alpha)g_2(T).$$

The problem is to show that  $g(T) \in V_{\delta}(T)$ . It is clear that g is linear and that

Therefore g is a positive linear functional and since

$$g(I) = \alpha g_1(I) + (1 - \alpha)g_2(I) = 1$$
, and  $||g|| = 1$ ,

g is in fact a state on  $\mathfrak U$ . Hence  $g(T) \in V_{\delta}(T)$ .

Having proved that both sets  $W_{\delta}(T)$  and  $V_{\delta}(T)$  are convex we shall now show that those two sets are actually equal.

THEOREM 1.4. For any  $T \in \mathfrak{L}(\mathfrak{S})$ ,  $W_{\delta}(T) = V_{\delta}(T)$ .

PROOF: It is easy to show that  $W_{\delta}(T) \subseteq V_{\delta}(T)$ . To prove the converse, we shall assume that  $\lambda \in V_{\delta}(T)$  and  $\lambda \notin W(T)$  and deduce a contradiction. Since  $\lambda \in V_{\delta}(T)$ , it follows that there exists a state f in  $\mathfrak{U}^*$  such that  $f(T) = \lambda$  and  $f(T^*T) \geqslant \delta^2$ .  $W_{\delta}(T)$  is convex. By rotating T, we may assume that

$$\Re e W_{\delta}(T) \leqslant \Re e \lambda - \alpha, \quad \alpha > 0.$$

Let

$$G = \{x \in \mathfrak{S}, ||x|| = 1 \text{ and } \Re e \langle Tx, x \rangle \geqslant \Re e \lambda - \alpha/2, \alpha > 0\}$$

and

$$\eta = \sup\{\|Tx\| : x \in G\}.$$

Then  $\eta < \delta$ . The set G is nonempty because if it is not, then for all  $x \in \mathfrak{S}$ , ||x|| = 1, we shall have

$$\Re e \langle Tx, x \rangle < \Re e \lambda - \alpha/2, \quad \alpha > 0.$$

But since f is a weak\*-limit of convex combinations of vector states

$$\forall \varepsilon > 0 \,\exists \, N = N(\varepsilon) \colon \forall n > N, \quad |f_n(T) - f(T)| < \varepsilon.$$

Also we can find  $M = M(\varepsilon)$  such that for all n > M the following inequality will hold:

$$|f_n(T^*T) - f(T^*T)| < \varepsilon.$$

Take  $\varepsilon < \min\{\alpha/2, (\delta^2 - 2\delta\eta)/2\}$  and  $n > \max(N, M)$ . Since

$$f_n(T) = \sum_{i=1}^n \alpha_i \omega_{z_i}(T) = \sum_{i=1}^n \alpha_i \langle Tx_i, x_i \rangle$$

for  $0 \le \alpha_i \le 1$ , and  $\sum_{i=1}^n \alpha_i = 1$ , we have

$$egin{aligned} \Re e \, f_n(x) &= \Re e \, \sum_{i=1}^n lpha_i \omega_{x_i}(T) = \Re e \, \sum_{i=1}^n lpha_i \langle Tx_i, \, x_i 
angle \ &= \sum_{i=1}^n lpha_i \Re e \, \langle Tx_i, \, x_i 
angle \, < \Re e \, \lambda - lpha/2. \end{aligned}$$

But  $f_n(x) > f(T) - \varepsilon$  and therefore we see that  $\Re e f_n(x) > \Re e \lambda - \varepsilon$ , implying that  $\varepsilon > \alpha/2$ , which is a contradiction.

Now for all  $x_i \in G$ , we have  $||Tx_i|| \leq \eta$ . Since also  $f(T^*T) < f_n(T^*T) + \varepsilon$  and  $\delta^2 \leq f(T^*T)$  we obtain

$$\delta^{2} \leqslant f(T^{*}T) < f_{n}(T^{*}T) + \varepsilon = \sum_{i=1}^{n} \alpha_{i} ||Tx_{i}||^{2} + \varepsilon < \eta^{2} + (\delta^{2} - 2\delta\eta)/2$$

$$= \eta^{2} + 1/2(\delta - \eta)^{2} - \eta^{2}/2$$

$$= \eta^{2}/2 + 1/2(\delta - \eta)^{2} < \delta^{2},$$

which is a contradiction. So we see that  $\lambda \notin W_{\delta}(T)$  implying that  $\lambda \notin V_{\delta}(T)$ . Hence  $\lambda \in V_{\delta}(T)$  implies that  $\lambda \in W_{\delta}(T)$ , and so  $V_{\delta}(T) \subseteq W_{\delta}(T)$ .

## REFERENCES

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